Relational algebra: a Kleene algebra central to the mathematics of program construction

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II Jornadas Luso-Galegas de Álgebra e Computação
Braga, 10-10-2008

On maths and computing

Interaction between maths and computing:

- computers helping maths: theorem proving, computational maths etc
- maths helping computing: many examples, among which the algebra of programming (**AoP**)

While the former are widely acknowledged, among the latter **AoP** is known only to the initiated.

- This talk aims at framing **AoP** in its proper algebraic context while showing its relevance to program construction.

It all starts from semirings of computations [3]...

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Semirings of computations

Abstract notion of a computation:

Semiring $(S, +, \cdot, 0, 1)$ *inhabited by computations (eg. instructions, statements) where*

- $x \cdot y$ *(usually abbreviated to* xy*) captures* **sequencing**
- $x + y$ *captures* **choice** *(alternation)*
- 0 *means* **death**
- 1 *means* **skip** *(do nothing)*

Technically:

- $(S, \cdot, 1)$ is a monoid
- $(S, +, 0)$ is a Abelian monoid
- $(\cdot)$ distributes over $(+)$
- 0 annihilates $(\cdot)$

Idempotency

- If $x + x = x$ holds for all x, then

$$x \leq y \stackrel{\mathrm{def}}{=} x + y = y \tag{1}$$

  is a partial order.

- Clearly, $0 \leq x$ for all x and $(+)$ is the *lub* with respect to $\leq$:

$$x + y \leq z \iff x \leq z \wedge y \leq z \tag{2}$$

  NB: $z := x + y$ in (2) means $x + y$ is upper bound; $\Leftarrow$ means it is the **least** upper bound (*lub*).

Kleene algebras

A Kleene algebra [5] adds to semiring $(S, +, \cdot, 0, 1)$ the *Kleene star* operator $(^*)$ such that

$$y + x(x^* y) \leq x^* y \quad (3)$$

$$y + (y x^*) x \leq y x^* \quad (4)$$

and

$$y + xz \leq z \Rightarrow x^* y \leq z \quad (5)$$

$$y + zx \leq z \Rightarrow y x^* \leq z \quad (6)$$

These basically establish $x^* y$ and $y x^*$ as prefix points of (monotonic) functions $(y + x \cdot \_)$ and $(y + \_ \cdot x)$, respectively.

KATs (tests and domains)

KAT = Kleene algebra with tests

- every p below 1 ($p \leq 1$) is a **test** and such that, for every such p there is $\neg p$ (the *complement* of p) such that

$$p + \neg p = 1$$
$$p \cdot \neg p = 0 = \neg p \cdot p$$

- Recent addition to semirings (inc. KATs) of a *domain* operator $d(x)$ capturing "enabledness" and satisfying axioms

$$\begin{array}{rcl} d(x) & \leq & 1 \\ d(0) & = & 0 \\ d(x+y) & = & d(x) + d(y) \\ d(xy) & = & d(x\ d(y)) \\ x & \leq & d(x)x \end{array}$$

Binary relations

The algebra of **binary relations** is a well known KAT:

KAT	Binary relations	Description
$x \cdot y$	$R \cdot S$	composition
$x + y$	$R \cup S$	union
0	$\bot$	empty relation
1	id	identity relation
$x \leq y$	$R \subseteq S$	inclusion
$p, \neg p$	$R \subseteq id$, $\neg R = id - R$	coreflexive relations
$d(x)$	$\delta\, R$	domain of R

Moreover, they form a complete, distributive lattice once *glb*s

$$X \;\subseteq\; R \cap S \quad\Leftrightarrow\quad (X \subseteq R) \wedge (X \subseteq S) \tag{7}$$

and supremum $\top$ are added.

How useful are binary relations?

- Not much if regarded merely as "sets of pairs"
- Very useful indeed — as a device for the algebraization of logic — if regarded as "**arrows**" ie. morphisms of a particular allegory [4]
- Arrows bring about a **type discipline** which leads to good things such as parametric **polymorphism**, etc etc

Relations as morphisms

Binary relations are typed:

Arrow notation

Arrow $A \xrightarrow{R} B$ denotes a binary relation from A (source) to B (target).

A, B are types. Writing $B \xleftarrow{R} A$ means the same as $A \xrightarrow{R} B$.

Infix notation

The usual infix notation used in natural language — eg.

$$\texttt{John}\ \mathit{IsFatherOf}\ \texttt{Mary}$$

— and in maths — eg.

$$0 \leq \pi$$

— extends to arbitrary $B \xleftarrow{R} A$: we write

$$b\ R\ a$$

to denote that $(b, a) \in R$.

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Functions are relations

- Lowercase letters (or identifiers starting by one such letter) will denote special relations known as **functions**, eg. f, g, suc, etc.
- We regard **function** $f : A \longrightarrow B$ as the binary **relation** which relates b to a iff $b = f\ a$. So,

$$b\ f\ a \text{ literally means } b = f\ a$$

- Therefore, we generalize

$$\boxed{\begin{array}{c} B \xleftarrow{f} A \\ b = f\ a \end{array}} \quad \text{to} \quad \boxed{\begin{array}{c} B \xleftarrow{R} A \\ b\ R\ a \end{array}}$$

- So, **function** id is the equality (equivalence) **relation**:

$$b\ id\ a \text{ means the same as } b = a$$

Composition

Function composition

$$\begin{tikzcd} B & A \arrow[l, "f"'] & C \arrow[l, "g"'] \arrow[ll, bend left, "f \cdot g"] \end{tikzcd} \tag{8}$$

$$b = f(g\ c)$$

extends to $R \cdot S$ in the obvious way:

$$b(R \cdot S)c \quad \Leftrightarrow \quad \langle \exists\ a\ ::\ b\ R\ a \wedge a\ S\ c \rangle \tag{9}$$

Note how this rule *removes* quantifier $\exists$ when applied from right to left.

Converses

Every relation $B \xleftarrow{R} A$ has a **converse** $B \xrightarrow{R^\circ} A$ which is such that, for all a, b,

$$a(R^\circ)b \quad \Leftrightarrow \quad b\ R\ a \tag{10}$$

Note that converse commutes with composition

$$(R \cdot S)^\circ = S^\circ \cdot R^\circ \tag{11}$$

and cancels itself

$$(R^\circ)^\circ = R \tag{12}$$

Function converses

Function converses f°, g° etc. always exist (as **relations**) and enjoy the following (very useful) property:

$$(f\ b)R(g\ a) \quad \Leftrightarrow \quad b(f^\circ \cdot R \cdot g)a \tag{13}$$

cf. diagram:

$$\begin{tikzcd}
C & D \arrow[l, "R"'] \\
B \arrow[u, "f"] & A \arrow[l, "f^\circ \cdot R \cdot g"] \arrow[u, "g"']
\end{tikzcd}$$

Why *id* (really) matters

Terminology:

- Say R *is* *reflexive* iff $id \subseteq R$
  pointwise: $\langle \forall\ a\ ::\ a\ R\ a \rangle$
- Say R *is* *coreflexive* iff $R \subseteq id$
  pointwise: $\langle \forall\ b, a\ :\ b\ R\ a :\ b = a \rangle$

Define, for $B \xleftarrow{R} A$:

Kernel of R	**Image** of R
$A \xleftarrow{\ker R} A$ $\ker R \triangleq R^\circ \cdot R$	$B \xleftarrow{\mathrm{img}\, R} B$ $\mathrm{img}\, R \triangleq R \cdot R^\circ$

Example

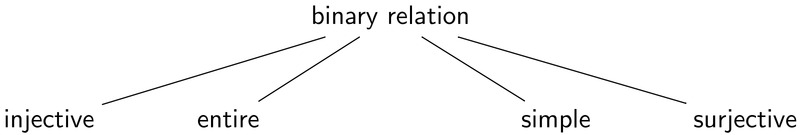
Kernels of functions:

$$
\begin{array}{ll}
 & a'(\ker f)a \\
\Leftrightarrow & \quad \{ \text{ substitution } \} \\
 & a'(f^{\circ} \cdot f)a \\
\Leftrightarrow & \quad \{ \text{ PF-transform rule (13) } \} \\
 & (f\ a') = (f\ a)
\end{array}
$$

In words: $a'(\ker f)a$ means a' *and* a *"have the same* f*-image"*

Binary relation taxonomy

Topmost criteria:



Definitions:

	Reflexive	*Coreflexive*
$\ker R$	entire R	injective R
$\mathrm{img}\, R$	surjective R	simple R

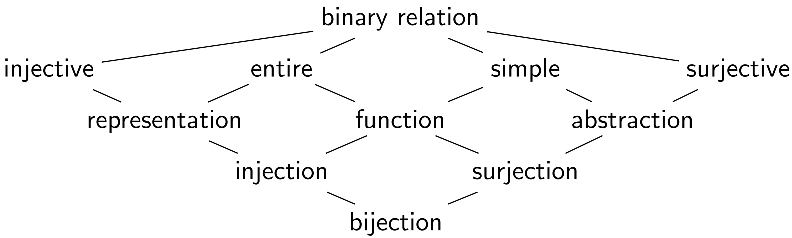
(14)

Facts:

$$\ker (R^\circ) = \mathrm{img}\, R \tag{15}$$
$$\mathrm{img}\, (R^\circ) = \ker R \tag{16}$$

Binary relation taxonomy

The whole picture:



(17)

Clearly:

- converse of *injective* is *simple* (and vice-versa)
- converse of *entire* is *surjective* (and vice-versa)
- smaller than injective (simple) is injective (simple)
- larger than entire (surjective) is entire (surjective)

Functions in one slide

A function f is a binary relation such that

Pointwise	**Pointfree**	
"Left" Uniqueness		
$b\ f\ a \wedge b'\ f\ a \ \Rightarrow\ b = b'$	$\mathrm{img}\, f \ \subseteq\ id$	(f is simple)
Leibniz principle		
$a = a' \ \Rightarrow\ f\ a = f\ a'$	$id \ \subseteq\ \ker f$	(f is entire)

which both together are equivalent to any of "al-gabr" rules

$$f \cdot R \subseteq S \ \Leftrightarrow\ R \subseteq f^{\circ} \cdot S \tag{18}$$

$$R \cdot f^{\circ} \subseteq S \ \Leftrightarrow\ R \subseteq S \cdot f \tag{19}$$

"Al-gabr" rules?

Recall *calculus of al-gabr and al-muqâbala* [1]:

al-gabr

$$x - z \leq y \quad \Leftrightarrow \quad x \leq y + z$$

al-hatt

$$x * z \leq y \quad \Leftrightarrow \quad x \leq y * z^{-1} \qquad (z > 0)$$

al-muqâbala

Ex:

$$4x^2 + 3 = 2x^2 + 2x + 6 \quad \Leftrightarrow \quad 2x^2 = 2x + 3$$

[1] Cf. *Kitâb al-muhtasar fi hisab al-gabr wa-almuqâbala* by Abû Al-Huwârizmî, the famous 9c Persian mathematician.

Example: function equality

Equating functions means comparing them in either way:

$$f = g \quad \Leftrightarrow \quad f \subseteq g \quad \Leftrightarrow \quad g \subseteq f \tag{20}$$

Calculation:

$$\begin{array}{ll}
 & f \subseteq g \\
\Leftrightarrow & \qquad \{ \text{ “al-gabr” (18) on } f \ \} \\
 & id \subseteq f^{\circ} \cdot g \\
\Leftrightarrow & \qquad \{ \text{ “al-gabr” (19) on } g \ \} \\
 & g^{\circ} \subseteq f^{\circ} \\
\Leftrightarrow & \qquad \{ \text{ converses } \} \\
 & g \subseteq f
\end{array}$$

A "Laplace transform analog" for logical quantification

The pointfree (PF) transform

ϕ	$PF\ \phi$
$\langle \exists\ a\ ::\ b\ R\ a \wedge a\ S\ c \rangle$	$b(R \cdot S)c$
$\langle \forall\ a, b\ ::\ b\ R\ a \Rightarrow b\ S\ a \rangle$	$R \subseteq S$
$\langle \forall\ a\ ::\ a\ R\ a \rangle$	$id \subseteq R$
$\langle \forall\ x\ ::\ x\ R\ b \Rightarrow x\ S\ a \rangle$	$b(R \setminus S)a$
$\langle \forall\ c\ ::\ b\ R\ c \Rightarrow a\ S\ c \rangle$	$a(S / R)b$
$b\ R\ a \wedge c\ S\ a$	$(b, c)\langle R, S \rangle a$
$b\ R\ a \wedge d\ S\ c$	$(b, d)(R \times S)(a, c)$
$b\ R\ a \wedge b\ S\ a$	$b\ (R \cap S)\ a$
$b\ R\ a \vee b\ S\ a$	$b\ (R \cup S)\ a$
$(f\ b)\ R\ (g\ a)$	$b(f^\circ \cdot R \cdot g)a$
TRUE	$b\ \top\ a$
FALSE	$b\ \bot\ a$

What do $\langle R, S \rangle$, $R \times S$ etc mean?

Forks for tupling

The **fork** ("split") combinator is essential for transforming predicates holding more than two quantified variables. From the definition,

$$(b, c)\langle R, S\rangle a \quad \Leftrightarrow \quad b\ R\ a \wedge c\ S\ a$$

which PF-transforms to

$$\langle R, S\rangle \quad = \quad \pi_1^\circ \cdot R \cap \pi_2^\circ \cdot S \tag{21}$$

we infer diagram

$$\begin{tikzcd}
A & A \times B \arrow[l, "\pi_1"'] \arrow[r, "\pi_2"] & B \\
& C \arrow[ul, "R"] \arrow[u, "\langle R,S\rangle"] \arrow[ur, "S"'] &
\end{tikzcd}$$

and "al-gabr" rule (Galois connection)

$$\pi_1 \cdot X \subseteq R \ \wedge \ \pi_2 \cdot X \subseteq S \quad \Leftrightarrow \quad X \subseteq \langle R, S\rangle \tag{22}$$

Coproducts for *"if-then-else'ing"*

Define dual ("either") combinator as

$$[R\ ,S] \;=\; (R \cdot i_1^\circ) \cup (S \cdot i_2^\circ) \quad \text{cf.} \quad \begin{tikzcd} A \arrow[r, "i_1"] \arrow[dr, "R"'] & A + B \arrow[d, "{[R\ ,S]}"] & B \arrow[l, "i_2"'] \arrow[dl, "S"] \\ & C & \end{tikzcd}$$

From this and the *lub* rule (2) we infer another "al-gabr" rule (Galois connection)

$$[R\ ,S] \subseteq X \quad \Leftrightarrow \quad R \subseteq X \cdot i_1 \;\wedge\; S \subseteq X \cdot i_2 \tag{23}$$

In fact, the stronger universal property holds:

$$[R\ ,S] = X \quad \Leftrightarrow \quad R = X \cdot i_1 \;\wedge\; S = X \cdot i_2 \tag{24}$$

Multiplying and adding relations

From "fork" and "either" derive

$$R \times S \triangleq \langle R \cdot \pi_1, S \cdot \pi_2 \rangle \tag{25}$$

$$R + S = [i_1 \cdot R \;, i_2 \cdot S] \tag{26}$$

whose pointwise meaning is, as given earlier:

ϕ	$PF\ \phi$
$a\ R\ c \wedge b\ S\ c$	$(a, b)\langle R, S\rangle c$
$b\ R\ a \wedge d\ S\ c$	$(b, d)(R \times S)(a, c)$

Absorption properties:

$$\langle R \cdot X, S \cdot Y \rangle = (R \times S) \cdot \langle X, Y \rangle \tag{27}$$

$$[R \;, S] \cdot (X + Y) = [R \cdot X \;, S \cdot Y] \tag{28}$$

$+$ meets $\times$

From both (22) and (24) we easily infer the **exchange law**,

$$[\langle R, S\rangle \ , \langle T, V\rangle] \quad = \quad \langle [R \ , T], [S \ , V]\rangle \tag{29}$$

holding for all relations as in diagram

$$\begin{tikzcd}
A \arrow[r, "i_1"] \arrow[d, "R"'] \arrow[drr, "S"', near end] & A + B & B \arrow[l, "i_2"'] \arrow[d, "V"] \arrow[dll, "T", near start] \\
C & C \times D \arrow[l, "\pi_1"] \arrow[r, "\pi_2"'] & D
\end{tikzcd}$$

Inductive relations

Example — inductive definition of $\geq$ over the natural numbers: for all $y, x \in \mathbb{N}_0$, define $\mathbb{N}_0 \xleftarrow{\geq} \mathbb{N}_0$ as the **least** relation satisfying

$$\begin{array}{l} y \geq 0 \\ y \geq x \ \Rightarrow \ (y+1) \geq (x+1) \end{array}$$

Thanks to (13), these clauses PF-transform to

$$\begin{array}{rcl} \top & \subseteq & \geq \cdot \underline{0} \\ \geq & \subseteq & suc^{\circ} \cdot \geq \cdot suc \end{array}$$

where $\underline{0}$ denotes the everywhere 0 constant function.

Least prefix points

We reason:

$$\begin{array}{l}
\left\{\begin{array}{lcl} \top & \subseteq & \geq \cdot \underline{0} \\ \geq & \subseteq & suc^\circ \cdot \geq \cdot suc \end{array}\right. \\
\Leftrightarrow \qquad \{ \text{ al-gabr (18) ; coproducts } \} \\
[\top\ , suc \cdot \geq] \ \subseteq\ \geq \cdot [\underline{0}\ , suc] \\
\Leftrightarrow \qquad \{ \text{ “al-gabr” (19) } \} \\
[\top\ , suc \cdot \geq] \cdot [\underline{0}\ , suc]^\circ \ \subseteq\ \geq \\
\Leftrightarrow \qquad \{ \text{ absorption property (28) } \} \\
[\top\ , suc] \cdot (id + \geq) \cdot [\underline{0}\ , suc]^\circ \ \subseteq\ \geq
\end{array}$$

In summary: $\geq$ is the least **prefix** point of monotonic function

$$f\ X \ \triangleq\ [\top\ , suc] \cdot (id + X) \cdot [\underline{0}\ , suc]^\circ$$

Diagrams help

Recognizing $[\underline{0}\ , suc] = in$ as initial $(1 + \_)$-algebra with carrier N_0 (Peano isomorphism) we draw

$$\begin{tikzcd}[column sep=huge, row sep=large]
N_0 \arrow[r, bend left, "in^\circ"] \arrow[r, phantom, "\cong"] \arrow[d, "\geq"'] & 1 + N_0 \arrow[l, bend left, "in"] \arrow[d, "id+\geq"] \arrow[dl, phantom, "\supseteq"] \\
N_0 & 1 + N_0 \arrow[l, "{[\top\ , suc]}"]
\end{tikzcd}$$

$$[\top\ , suc] \cdot (id + \geq)\ \subseteq\ \geq \cdot in$$

Since $[\top\ , suc]$ uniquely determines $\geq$ (least prefix points are unique, etc), we resort to the popular notation

$$\geq\ =\ (\!|[\top\ , suc]|\!) \qquad (30)$$

to express this fact. (See summary of general theory in the sequel.)

Introducing the $\kappa\alpha\tau\alpha$ combinator

In general, for F a polynomial functor (relator) and
$\mu\mathsf{F} \xleftarrow{in} \mathsf{F}(\mu\mathsf{F})$ initial:

$$
\begin{tikzcd}
\mu\mathsf{F} \arrow[rr, bend left, "in^\circ"] \arrow[rr, phantom, "\cong"] \arrow[d, "(|R|)"'] & & \mathsf{F}(\mu\mathsf{F}) \arrow[ll, bend left, "in"] \arrow[d, "\mathsf{F}(|R|)"] \\
A & & \mathsf{F}\,A \arrow[ll, "R"]
\end{tikzcd}
$$

there is a unique solution to equation $X = R \cdot \mathsf{F}\,X \cdot in^\circ$ characterized by universal property:

$$X = (|R|) \;\Leftrightarrow\; X = R \cdot \mathsf{F}\,X \cdot in^\circ \qquad (31)$$

(Read $(|R|)$ as "$\kappa\alpha\tau\alpha$ R".)

Introducing the $\kappa\alpha\tau\alpha$ combinator

Therefore (cf. Knaster-Tarski) $(\!|R|\!)$ is both **the least** prefix point

$$(\!|R|\!) \subseteq X \quad \Leftarrow \quad R \cdot \mathsf{F}\, X \cdot in^{\circ} \subseteq X \tag{32}$$

and **the greatest** postfix point:

$$X \subseteq (\!|R|\!) \quad \Leftarrow \quad X \subseteq R \cdot \mathsf{F}\, X \cdot in^{\circ} \tag{33}$$

Corollaries include **reflexion**,

$$(\!|in|\!) \quad = \quad id \tag{34}$$

$\kappa\alpha\tau\alpha$-**fusion**,

$$S \cdot (\!|R|\!) \subseteq (\!|X|\!) \quad \Leftarrow \quad S \cdot R \subseteq X \cdot \mathsf{F}\, S \tag{35}$$

monotonicity,

$$(\!|R|\!) \subseteq (\!|X|\!) \quad \Leftarrow \quad R \subseteq X \tag{36}$$

etc.

Why $\kappa\alpha\tau\alpha$s?

- What's the advantage of writing $\geq \; = (\![\top \; , suc]\!)$? Is it just a matter of *style* or *economy* of notation?
- No: think of proving that $\geq$ is **transitive**:

$$\langle \forall \; x, y, z \; :: \; x \geq y \wedge y \geq z \Rightarrow x \geq z \rangle$$

Instead of providing an explicit (inductive) proof, we go *pointfree* and write:

$$\geq \cdot \geq \quad \subseteq \quad \geq$$

which instantiates $\kappa\alpha\tau\alpha$-**fusion** (35), for $R, X := [\top \; , suc]$.

Thank you, $\kappa\alpha\tau\alpha$-fusion

We reason:

$$
\begin{array}{cl}
 & \geq \cdot \geq \;\subseteq\; \geq \\
\Leftrightarrow & \quad \{ \text{ definition (30) } \} \\
 & \geq \cdot (\![\top , suc]\!) \;\subseteq\; (\![\top , suc]\!) \\
\Leftarrow & \quad \{ \; \kappa\alpha\tau\alpha\text{-fusion (35) } \} \\
 & \geq \cdot [\top , suc] \subseteq [\top , suc] \cdot (id + \geq) \\
\Leftrightarrow & \quad \{ \text{ coproducts (28, etc) } \} \\
 & \geq \cdot \top \subseteq \top \;\wedge\; \geq \cdot\, suc \subseteq suc \cdot \geq \\
\Leftrightarrow & \quad \{ \text{ everything is at most } \top \; \} \\
 & \geq \cdot\, suc \subseteq suc \cdot \geq \\
\Leftarrow & \quad \{ \; \geq \cdot\, suc = suc \cdot \geq \text{ (31) } \} \\
 & \textsc{True}
\end{array}
$$

By the way

Direct use of universal property (31) would lead to

$$
\begin{array}{ll}
 & \geq \; = \; (\![\top \; , suc]\!) \\
\Leftrightarrow & \qquad \{ \;\; (31) \;\; \} \\
 & \geq \cdot [\underline{0} \; , suc] = [\top \; , suc] \cdot (id + \geq) \\
\Leftrightarrow & \qquad \{ \;\; \text{expand, go pointwise, simplify} \;\; \} \\
 & \left\{ \begin{array}{l} y \geq 0 \\ y \geq (x+1) \Leftrightarrow y > 0 \wedge (y-1) \geq x \end{array} \right.
\end{array}
$$

So, the above and our starting (co-inductively flavored) definition

$$
\begin{array}{l}
y \geq 0 \\
y \geq x \;\Rightarrow\; (y+1) \geq (x+1)
\end{array}
$$

are *equivalent* (by construction).

$\kappa\alpha\tau\alpha$ meets fork

What about $\kappa\alpha\tau\alpha$s which are forks? We reason:

$$
\begin{array}{ll}
 & (|\langle R, S\rangle|) \subseteq \langle X, Y\rangle \\
\Leftarrow & \{ \text{ least prefix point (32) } \} \\
 & \langle R, S\rangle \cdot \mathsf{F}\langle X, Y\rangle \cdot in^{\circ} \subseteq \langle X, Y\rangle \\
\Leftrightarrow & \{ \text{ ``al-gabr'' rule (22) } \} \\
 & \left\{ \begin{array}{l} \pi_1 \cdot \langle R, S\rangle \cdot \mathsf{F}\langle X, Y\rangle \cdot in^{\circ} \subseteq X \\ \pi_2 \cdot \langle R, S\rangle \cdot \mathsf{F}\langle X, Y\rangle \cdot in^{\circ} \subseteq Y \end{array} \right. \\
\Leftarrow & \{ \; X := \langle R, S\rangle \text{ in (22); monotonicity } \} \\
 & \left\{ \begin{array}{l} R \cdot \mathsf{F}\langle X, Y\rangle \cdot in^{\circ} \subseteq X \\ S \cdot \mathsf{F}\langle X, Y\rangle \cdot in^{\circ} \subseteq Y \end{array} \right.
\end{array}
$$

Handling mutually recursive relations

- Rule

$$(\!|\langle R, S\rangle|\!) \subseteq \langle X, Y\rangle \quad \Leftarrow \quad \left\{ \begin{array}{l} R \cdot \mathsf{F}\langle X, Y\rangle \cdot in^\circ \subseteq X \\ S \cdot \mathsf{F}\langle X, Y\rangle \cdot in^\circ \subseteq Y \end{array} \right. \tag{37}$$

  tells us how to combine two mutually recursive relations into a single one.

- In the case of functions (20) we get equivalence

$$\left\{ \begin{array}{l} x \cdot in = r \cdot \mathsf{F}\langle x, y\rangle \\ y \cdot in = s \cdot \mathsf{F}\langle x, y\rangle \end{array} \right. \quad \Leftrightarrow \quad \langle x, y\rangle = (\!|\langle r, s\rangle|\!) \tag{38}$$

  known as "Fokkinga's mutual recursion theorem" [2].

- Both (37,38) generalize to $n > 2$ mutually recursive relations (functions) and can be used for program optimization.

Handling mutually recursive relations

- Notice that in° plays no special role in the calculation of (37); so it can be replaced by arbitrary (suitably typed) D.
- This generalizes rule (37) to **divide-and-conquer** algorithms described by recursive relations which are fixpoints of $f\ X \triangleq\ R \cdot (\mathsf{F}\,X) \cdot D$, where R describes the **conquer** step and D the **divide** step.
  (Btw, these are known as *hylomorphisms* [2].)
- For economy of presentation, the example which follows is a direct application of the special case where all relations are functions (38).

Example — exponential function

Taylor series:

$$e^x = \sum_{i=0}^{\infty} \frac{x^i}{i!} \tag{39}$$

Computing finite approximation (n terms)

$$e^x\ n = \sum_{i=0}^{n} \frac{x^i}{i!} \tag{40}$$

takes quadratic time. Wishing to calculate a linear-time algorithm from this mathematical definition, we first head for an inductive definition:

$$\begin{aligned} e^x\ 0 &= 1 \\ e^x\ (n+1) &= \underbrace{\frac{x^{n+1}}{(n+1)!}}_{h_x\, n} + \underbrace{\sum_{i=0}^{n} \frac{x^i}{i!}}_{e^x\ n} \end{aligned}$$

Example — exponential function

We thus get primitive recursive definition

$$\begin{aligned} e^x\ 0 &= 1 \\ e^x\ (n+1) &= h_x n + e^x\ n \end{aligned}$$

where $h_x n$ unfolds to $\frac{x^{n+1}}{(n+1)!} = \frac{x}{n+1}\frac{x^n}{n!}$. Therefore:

$$\begin{aligned} h_x 0 &= x \\ h_x(n+1) &= \frac{x}{n+2}(h_x n) \end{aligned}$$

Introducing $s2\ n = n+2$, we derive:

$$\begin{aligned} & s2\ 0 = 2 \\ & s2(n+1) = 1 + s2\ n \end{aligned}$$

Example — exponential function

We can thus put e^x, $s2$ and h_x together in a system of three mutually recursive functions e^x, $s2_x$ and h_x over the naturals, which PF-transform to

$$
\begin{aligned}
e^x \cdot in &= \underbrace{[\underline{1}\ , (+) \cdot \langle \pi_1, \pi_2 \cdot \pi_2 \rangle]}_{r} \cdot \mathsf{F}\langle e^x, \langle s2_x, h_x \rangle\rangle \\
s2_x \cdot in &= \underbrace{[\underline{2}\ , suc \cdot \pi_1 \cdot \pi_2]}_{s} \cdot \mathsf{F}\langle e^x, \langle s2_x, h_x \rangle\rangle \\
h_x \cdot in &= \underbrace{[\underline{x}\ , (*) \cdot ((x/) \times id) \cdot \pi_2]}_{t} \cdot \mathsf{F}\langle e^x, \langle s2_x, h_x \rangle\rangle
\end{aligned}
$$

respectively, for

$$
\begin{aligned}
in &= [\underline{0}\ , suc] \\
\mathsf{F}\, X &= id + X
\end{aligned}
$$

Example — exponential function

From this system we obtain, thanks to the mutual recursion law (38)

$$\begin{array}{rcl} aux_x & \triangleq & \langle e^x, \langle s2_x, h_x \rangle\rangle \\ & = & \{\ (38)\ \} \\ & & (\!|\langle r, \langle s, t\rangle\rangle|\!) \end{array}$$

for

$$\begin{array}{rcl} r & = & [\underline{1}\ , (+) \cdot \langle \pi_1, \pi_2 \cdot \pi_2 \rangle] \\ s & = & [\underline{2}\ , suc \cdot \pi_1 \cdot \pi_2] \\ t & = & [\underline{x}\ , \underbrace{(*) \cdot ((x/) \times id) \cdot \pi_2}_{u}] \end{array}$$

Example — exponential function

Next we apply the exchange law (29) to $\langle r, \langle s, t \rangle \rangle$ (twice):

$$\langle r, \langle s, t \rangle \rangle \quad = \quad [\langle \underline{1}, \langle \underline{2}, \underline{x} \rangle \rangle \ , \langle (+) \cdot \langle \pi_1, \pi_2 \cdot \pi_2 \rangle, \langle suc \cdot \pi_1 \cdot \pi_2, u \rangle \rangle]$$

Thanks to universal properties (31) and (22) [2] we obtain

$$\begin{aligned} aux_x \cdot \underline{0} &= \langle \underline{1}, \langle \underline{2}, \underline{x} \rangle \rangle \\ aux_x \cdot suc &= \langle (+) \cdot \langle \pi_1, \pi_2 \cdot \pi_2 \rangle, \langle suc \cdot \pi_1 \cdot \pi_2, u \rangle \rangle \cdot aux_x \\ e^x &= \pi_1 \cdot aux_x \end{aligned}$$

that is, we have calculated linear implementation

[2] For functions.

Example — exponential function

```
exp x n = let (e,b,c) = aux x n
          in e where
             aux x 0 = (1,2,x)
             aux x (i+1) = let (e,s,h) = aux x i
                           in  (e+h,s+1,(x/s)*h)
```

which can be identified as the denotational semantics of a while loop, encoded below in the C programming language:

```
float exp(float x, int n)
{
  float e=1; int s=2; float h=x; int i;
  for (i=0;i<n+1;i++) {e=e+h;h=(x/s)*h;s++;}
  return e;
};
```

Summing up

- Algebra of Programming (**AoP**): calculating ("correct by construction") programs from specifications
- Pointfree notation: Tarski's *set theory without variables* [7]
- Kleene algebra of (typed) relations: arrows (not points) provide further structure while ensuring type checking
- *Ut faciant opus signa*:

> [Symbolisms] *"have invariably been introduced to make things easy. [...] by the aid of symbolism, we can make transitions in reasoning almost mechanically by the eye, which otherwise would call into play the higher faculties of the brain. [...] Civilisation advances by extending the number of important operations which can be performed without thinking about them."*
>
> *(Alfred Whitehead, 1911)*

However

Despite textbooks such as [2], **Algebra of Programming** is still land of nobody. Why?

- Software theorists: too busy with their pre-scientific theories (if any)
- Algebraists: not sufficiently aware of program construction as a mathematical discipline
- Both: the required background (categories, allegories, etc) is most often found missing from undergrad curricula.

Selected topic of interest

- Pointfree notations are emerging elsewhere in the context of eg. digital signal processing (SPIRAL project, CMU [6]) which abstract linear signal transforms in terms of (index-free) matrix operators.
- Kleene algebras scale up to the corresponding **matrix Kleene algebras** [1]
- Parallel with relational algebra is obvious.
- Following a similar path, we want to investigate the "matrices as arrows" approach purported by **categories of matrices** (PhD project).
- We believe a better (typed!) calculus of (Kleene) matrix algebras will emerge which will improve reasoning about linear transforms in DSP, divide-and-conquer algorithms, etc.

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